A geometric approach for three-phase load balancing in distribution networks

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Publication Details
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**Keywords**
networks, load, distribution, phase, three, approach, geometric, balancing

**Disciplines**
Engineering | Science and Technology Studies

**Publication Details**

This conference paper is available at Research Online: [http://ro.uow.edu.au/eispapers/4426](http://ro.uow.edu.au/eispapers/4426)
A GEOMETRIC APPROACH FOR THREE-PHASE LOAD BALANCING IN DISTRIBUTION NETWORKS

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Abstract: Network reconfiguration for load balancing is the determination of switching-options for a particular set of loads on a distribution system that relieves overloads and improves load balancing. It is performed by altering the topological structure of distribution feeders which is done by changing the open/closed status of sectionalizing and tie-switches. To relieve overloads of a network and to balance feeder loads, the loads are transferred from the heavily loaded feeders to the lightly loaded ones. In this paper, a geometric approach for load balancing is presented. In the proposed method, each loop in a network is represented as a circle, which is derived from the relationship between the change of load balancing due to the branch-exchange and the power-flows in the branches. If there is no change of load balancing in the system, then all the circles touch each other at the (0,0) coordinate. The circles with no load balancing are called zero load balancing change circles. The maximum load balancing loop in the network is identified by comparing the radii of all the modified zero load balancing change circles. The corresponding loop of the largest one gives the maximum improvement of load balancing in the network. Then the possible branch exchanges in the maximum load balancing loop are investigated by comparing the size of the circle for every branch-exchange. If the loads are balanced due to a branch exchange, the size of the circle diminishes and hence the smallest circle gives the maximum improvement in load balancing. To show the efficiency and performance of the proposed method for the solution of computationally complex and large dimensionality problems, a system with 69-bus and 7 laterals has been considered as the test system. Test results have shown that the proposed method can identify the most effective branch-exchange operations for load balancing with less computational effort and time. The number of load flow solutions has been reduced to a greater extent in the proposed method.

Keywords: Distribution networks, Load balancing, Geometric, and Network reconfiguration.

I. INTRODUCTION

The increase in power demand and high load density in the urban areas makes the operation of distribution systems more complex to control. To meet the load demand, the requirements are met by increasing the substation capacity and the number of feeders. Due to financial and other constraints, the requirements could not be met. Therefore, load-balancing techniques are resorted to provide more capacity margin for the substation to meet the extra load demand. With the advent of fast computing systems, the advanced control of electrical power systems has become a viable one. For effective coordination of their protective systems, the distribution systems are normally configured radially. Most distribution networks use sectionalizing-switches that are normally closed, and tie-switches that are normally opened. These switches are used for both protection and network reconfiguration. Due to changing operating conditions, networks are reconfigured to reduce the system power loss i.e., network reconfiguration for loss reduction, and to relieve overloads in the networks i.e., network reconfiguration for load balancing. By modifying the radial structure of the feeders periodically, they significantly improve the operating conditions of the overall system. Feeders in a distribution system normally have a mixture of industrial, commercial, residential, lighting etc. type of loads. The peak load on the substation transformers and feeders occurs at different times of the day. Therefore, the distribution system becomes heavily loaded at certain times of the day, and lightly loaded at other times. Load balancing is obtained by transferring loads from the heavily loaded feeders to lightly loaded feeders by reconfiguring the network. This is done in order to reschedule the loads more efficiently for maximizing the load balancing in the system. Reconfiguration also allows smoothing out the peak demands, improving the voltage profile in the feeders and increasing the network reliability.

Aoki et. al. [1,2] have formulated the load balancing and service restoration problems by considering the capacity and voltage constraints as a mixed integer non-linear optimization problem. Baran et. al. [3] have formulated the problem of loss minimization and load balancing as an integer-programming problem. Chiang et. al. [4] have proposed a constrained multi-objective and non-
A new approach for network reconfiguration based load balancing is presented in this paper by introducing a geometrical technique. The method can be implemented by measuring and comparing the radii of the circles, which are formed from the relation of power-flows in the branches. The first stage of the method determines the maximum load-balancing loop graphically. In a distribution system a loop is associated by a tie-line and hence there are several loops in the system. To obtain the maximum load-balancing loop, the zero load balancing-change circles are drawn and the loop with the largest circle is identified for maximum load balancing. The second stage determines the switching-operation to be executed in that loop to reach a maximum improvement in load balancing by network reconfiguration. Lin et al. [8] have presented a new load balancing algorithm for the three-phase unbalanced distribution systems. The proposed algorithm of [8] could confine the size of the solution state-space, and avoid complex numerical computations. M. A. Kashem et al. [9] have developed a technique for load balancing by measuring the distances of various points representing power-flows.

Consider the branch exchange between branches t and m where t is the open branch and m is the closed branch in Fig. 1. The loop is divided into two parts by the open branch t in which the first part consisting the branches in the loop that extends between nodes 0 to l and the second part consisting the branches that extends between nodes 0 to h. The change in load balancing formula due to this branch exchange as given in reference 9 is rewritten as follows:

\[ \Delta LB_{sm}^S = 2P_m \sum k_i P_i - \sum k_i P_i + 2Q_m \sum k_i Q_i - \sum k_i Q_i - (P_m^2 + Q_m^2)k_{loop} \]  

where,

\[ \Delta LB_{sm}^S : \text{the change in load balancing due to branch exchange } t-m, \]

\[ P_m \text{ and } Q_m : \text{real and reactive power flows in branch } m \] (to be removed) respectively,

\[ k_i : \text{the inverse of the square of the maximum capacity of a branch } \]

\[ k_i P_i \text{ and } k_i Q_i : \text{products of } k P \text{ and } P \text{ on the lv-side } l \text{ and hv-side } h \text{ of the loop respectively,} \]

\[ k_{loop} : \text{the sum of } k \text{ value of each branch in the loop in which power-flow exists.} \]

If the value of \( \Delta LB_{sm}^S \) is greater than zero, the load balancing is improved, and if it is less than zero, the load balancing is deteriorated.

Equation (1) can be rearranged as shown below to be represented by a circle,

\[ \left[ \frac{P_m - A/C}{C} \right]^2 + \left[ \frac{Q_m - B/C}{C} \right]^2 = \left[ \frac{A/C}{C} \right]^2 - \Delta LB_{sm}^S \]  

where, A = \sum k_i P_i \text{ and } B = \sum k_i Q_i \text{ and } C = k_{loop}

The center of the circle is \( (A/C, B/C) \) and its radius is

\[ \left[ \frac{A^2 + B^2}{C^2} - \frac{\Delta LB_{sm}^S}{C} \right]^{1/2} \]

The zero load balancing change circle is the one which gives the change in load balancing to be zero (i.e. \( \Delta LB_{sm}^S = 0 \)). It can be represented as follows:

\[ \left[ \frac{P_m - A/C}{C} \right]^2 + \left[ \frac{Q_m - B/C}{C} \right]^2 = \left[ \frac{A/C}{C} \right]^2 + \left[ \frac{B/C}{C} \right]^2 \]  

Hence the radius of the circle is \( \left[ \frac{A^2 + B^2}{C^2} \right]^{1/2} \) for zero load balancing change (i.e. \( \Delta LB_{sm}^S = 0 \)) and is decreased if load balancing is improved in the loop.
II.A GEOMETRICAL DETERMINATION OF THE MAXIMUM LOAD BALANCING LOOP

The basic idea of the proposed search scheme is by first identifying the maximum load balancing loop in the network and then a switching-option is determined for the loop that gives a maximum improvement in load balancing.

Equation (2) can be rearranged as,

\[ \Delta L_{B_{sm}} = (A^2 + B^2)C - C((P_m - A/C)^2 + (Q_m - B/C)^2) \]  

(4)

From equation (4), it is evident that the change in load balancing, \( \Delta L_{B_{sm}} \) will be maximum when \( P_m = A/C \) and \( Q_m = B/C \). Therefore, the maximum value of \( \Delta L_{B_{sm}} \) in a loop is

\[ \Delta L_{B_{sm \, loop}} = (A^2 + B^2)/C \]

(5)

The equation for \( \Delta L_{B_{sm \, loop}} \) is similar to the radius of the zero load balancing change circle except the constant of multiplication in the equation.

Rearranging equation (3), we get the modified zero-load balancing change circle as,

\[ (P_m - A/C)^2 + (Q_m - B/C)^2 = \frac{A^2 + B^2}{C} \]

(6)

where, \( P_m = \sqrt{C} P_m \) and \( Q_m = \sqrt{C} Q_m \)

The radius of the above circle is \( \left[ \frac{A^2 + B^2}{C} \right]^{1/2} \) which is the square root of \( \Delta L_{B_{sm}} \) value. The loop circles can be drawn using equation (6) and the largest circle will give the maximum loss reduction loop among all the circles drawn for all the loops in the network.

The drawback of equation (5) is that the value of the \( \Delta L_{B_{sm \, loop}} \) is always a positive constant for a particular loop as it is independent of \( P_m \) and \( Q_m \). However, load balancing depends on the value of \( \Delta L_{B_{sm}} \) i.e. if it is +ve, load balancing is improved and if it is -ve, load balancing is deteriorated [9]. Whether the load balancing is improved or deteriorated cannot be shown by equation (5) because it is always positive. Therefore to ensure that the largest circle gives the maximum improvement in load balancing, the nominal branch is considered. The nominal branch is the first adjacent branch to the tie branch on the lower voltage side of the loop and the nominal load balancing is the load balancing which occurs by exchanging the open branch with the nominal branch. If the nominal load balancing is negative, then there is no branch in the loop that can be a candidate for branch exchange [9]. If the nominal load balancing is positive, it means that maximum load balancing is achieved and the branch in the loop for branch exchange is determined.

Equation (4) can be rewritten for nominal branch exchange as,

\[ \Delta L_{B_{sm}} = \left( A^2 + B^2 \right) C - C((P_k - A/C)^2 + (Q_k - B/C)^2) \]  

(7)

where, \( P_k \) and \( Q_k \) are the power flows in the nominal branch in the network.

Using equation (7) the nominal loop circle is drawn for nominal branch exchange and compared with the zero load balancing change circle drawn by using equation (3) in the respective loop. If the nominal loop circle is reduced in size, then the nominal load balancing is positive and the branch exchange in the loop gives the maximum load balancing. Otherwise the next largest circle is considered and checked as above.

II.B GEOMETRICAL DETERMINATION OF SWITCHING-OPTION FOR LOAD BALANCING

In network reconfiguration, multiple switching-options are possible for optimal or near optimal configuration where several tie and sectionalizing switches are simultaneously closed and/or opened by the successive application of the proposed scheme. The best switching-option to be implemented is chosen in each successive operation that maximizes load balancing the most, without violating the constraints such as voltage, capacity of lines and transformers and reliability.

The authors in [9] have established a rule for the switching exchange of switches by considering the lower and higher voltage side of the loop. In ref. 9, load balancing is achieved when loads are transferred from the lower voltage side of a loop to the higher voltage side of that loop. It was noted that each line contains a sectionalizing switch to be operated for network reconfiguration. In this paper, branches of lower voltage side of a loop are considered for branch exchange and the most suitable switching-exchange is found by the knowledge of geometry as described below.

After determining the loop that would give the maximum load balancing by the geometrical procedure described in the previous section, a branch-exchange is identified in that loop. The various circles are drawn by using equation (4) for all the branches in the lower voltage side of the loop.

The size of the circle depends on the value of \( \Delta L_{B_{sm}} \). It gets reduced as the \( \Delta L_{B_{sm}} \) value is positive and increased if \( \Delta L_{B_{sm}} \) is negative. When \( \Delta L_{B_{sm}} \) is positive maximum for a branch-exchange, then the corresponding circle will be the smallest one. Therefore, the smallest circle will give the best solution and the corresponding branch-exchange will give the maximum load balancing in the network.

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III. SOLUTION TECHNIQUE

The procedure to determine the loop, which gives the maximum improvement in load balancing, is summarized below.

i) Run the load flow program to obtain the estimates of the power flows in the branches.

ii) Determine the radii of the modified zero-load balancing change circles for all the loops in the system (equation (6)) and compare. Select the largest circle and the corresponding loop.

iii) Check the nominal branch-exchange in the selected loop whether it contributes positive load balancing or not. This is done by comparing the radius of the nominal loop circle (equation (7)) with the radius of zero-load balancing change circle (equation (3)). If nominal loop circle is reduced in size, then go to step (v). Otherwise select the next largest circle and repeat step (iii).

iv) If all the nominal loop circles are increased in size compared to their respective zero-load balancing change circles, then stop.

v) Identify the loop of the corresponding circle and declare it as the maximum load balancing loop.

After finding the loop that gives the maximum load balancing, the next step is to determine the branch to be exchanged.

vi) Start from the nominal branch of the selected loop in step (v) and search the branches backward on the lower voltage side of the loop until the radius of the loop circle (equation (4)) is found to be minimum. Identify the corresponding branch-exchange.

vii) Check the selected branch-exchange for constraints violation. If all the constraints are satisfied, go to step (viii); otherwise discard the selected branch-exchange and continue the search forward and backward (one by one), until the constraints are satisfied and also the radius is found to be minimum.

viii) Declare the selected branch-exchange for maximum load balancing.

ix) Repeat steps (i) to (viii) using the selected branch-exchange.

IV. TEST RESULTS AND DISCUSSIONS

The test system is a hypothetical 12.66 kV radial distribution system with 69 buses and 7 laterals, which has been derived from a portion of the PG&E distribution system [7]. The schematic diagram of the test system is shown in Figure 2 and the relevant data are same as in ref. 7. There are 5 tie-lines (looping branches) in the system and sectionalizing switches on every branch of the system.

The total system loads for the base configuration are 3802.19 kW and 2694.60 kVAR.

The proposed geometrical method is tested on the 69-bus test system and test results are given in Table 1. It is shown that at the first search-level, at which, the radii of modified zero-load balancing change circles for all loops are calculated, and the loop that gives the largest value of radius is selected. In the next step, a comparison is made between the nominal loop circle and the zero-load balancing change loop circle to make sure that the selected loop gives the improvement in load balancing. The nominal circle has to be smaller to achieve load balancing in the system. In the next stage, the branch-exchange in that loop is selected by determining and comparing the radii of the loop circles starting from the nominal branch and searching the branches backward in the lower voltage side of the selected loop until the smallest circle is found to give the maximum improvement in load balancing. As shown in Table 1, the first iteration of the proposed method has found loop 4 as the maximum load balancing loop, and the branch exchange 73-59 to be performed in that loop to achieve maximum load balancing as shown in Figures 3 and 4 respectively. In Figure 3, the largest circle is found as the circle for the loop 4, which gives the maximum load-balancing loop. The loop circles are drawn for all the branch-exchanges in the lower voltage side of the loop 4 in Figure 4(a). It is visualized that the smallest circle is the circle for branch-exchange 73-59. In Figure 4(b) five small circles are shown and the smallest one is the circle for the...
branch-exchange 73-59, which gives the maximum load balancing. The above procedure is continued until all the loops have the nominal circles larger compared to their respective zero-load balancing change circles. It is to be noted that there is no need to draw the loop circles and the decision can be taken by comparing the radii of the circles only.

The proposed method considers the branches in the loop, which gives maximum improvement in load balancing by considering one branch-exchange in each search level. Because of lesser number of load flow solutions it is seen that the proposed method has greatly reduced computational time compared to Baran and Wu's method [3]. It is experimentally proven that the methods involved with a large number of load flow solutions, sometimes, seem exhaustive and unrealistic specially for larger systems, as they require more computational time.

Table 2 shows the summary of the selection of branch-exchanges for load balancing in all the search-levels. The optimal network configuration for load balancing (shown in Figure 5) is achieved after 5 search-levels of the proposed method, in which, at each search level (or iteration) a load flow solution is obtained, the maximum load balancing loop is selected and a branch-exchange is determined. In the fifth search-level, all the nominal circles become larger in size compared to their respective zero-load balancing change circles and these are shown in Table 2.

Table 1: Selection of Branch-Exchange in the First Search-Level of the Proposed Method

<table>
<thead>
<tr>
<th>Loop No.</th>
<th>Radius of mod. zero-load balancing loop circle</th>
<th>Selected loop</th>
<th>Radius of nominal loop circle</th>
<th>Radius of zero-load balancing change loop circle</th>
<th>Radius of smallest circle</th>
<th>Corresponding branch-exchange (in - out)</th>
</tr>
</thead>
<tbody>
<tr>
<td>1</td>
<td>0.0765</td>
<td>Loop 4</td>
<td>0.0075 (Smaller)</td>
<td>0.1725 (larger)</td>
<td>0.0060</td>
<td>73 - 59</td>
</tr>
<tr>
<td>2</td>
<td>0.0225</td>
<td></td>
<td></td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>3</td>
<td>0.0750</td>
<td></td>
<td></td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>4</td>
<td>0.1575</td>
<td></td>
<td></td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>5</td>
<td>0.0600</td>
<td></td>
<td></td>
<td></td>
<td></td>
<td></td>
</tr>
</tbody>
</table>

Note: Radius value is in per unit.

Table 2: Selection of Branch-Exchanges by the Proposed Geometrical Method

<table>
<thead>
<tr>
<th>Search Level</th>
<th>Loop Selection</th>
<th>Confirmation of load balancing</th>
<th>Identification of branch-exchange</th>
</tr>
</thead>
<tbody>
<tr>
<td>1</td>
<td>0.1575</td>
<td>4</td>
<td>0.0075 (Smaller) 0.1725 (larger) 0.0060 73 - 59</td>
</tr>
<tr>
<td>2</td>
<td>0.1275</td>
<td>3</td>
<td>0.0270 0.0900 0.1500 0.0975 74 - 62</td>
</tr>
<tr>
<td>3</td>
<td>0.2370</td>
<td>5</td>
<td>0.0900 0.1500 0.1500 0.0975 74 - 62</td>
</tr>
<tr>
<td>4</td>
<td>0.1425</td>
<td>3</td>
<td>0.0225 0.0975 0.0975 0.0975 13 - 15</td>
</tr>
<tr>
<td>5</td>
<td>0.4470</td>
<td>1</td>
<td>0.0750 0.0225 0.0225 0.0225 13 - 15</td>
</tr>
</tbody>
</table>

Note: Radius value in per unit.

'Change' indicates branch into the system and branch out from the system respectively.

* Loop 1 is associated with tie-line 70, loop 2 with tie-line 71, loop 3 with tie-line 72, loop 4 with tie-line 73 and loop 5 with tie-line 74.
V. CONCLUSION

A geometrical approach for load balancing has been presented in this paper. In the proposed method, various circles, which are formed from the power flows in the branches, are compared. A two-stage searching approach to the switching operation is proposed. In the first stage, the modified zero-load balancing circles are compared for various loops and the largest one is selected as the maximum load balancing loop. To ensure the load balancing for a branch-exchange in the selected loop, the nominal loop circle is compared with the zero-load balancing change circle. If the nominal loop circle-size is reduced, then the corresponding loop gives the maximum load balancing. Otherwise the next largest circle is considered. In the second stage, the circles for various branch-exchanges in the selected loop are compared and the smallest one is identified for the best solution. The corresponding branch-exchange of the smallest circle gives the maximum load balancing in the system. The proposed method has been applied to a 69-bus test system and the improvement in efficiency and the performance are established. From the test results it is found that the proposed method has identified the most effective branch-exchange operations for load balancing.

REFERENCES